

**Amendments to the Claims:**

The following listing of claims will replace all prior versions, and listings, of claims in the application:

1. (Previously Presented) A method for generating a one-way function dependent on a one-way function  $H$  and a unique value  $d$  for a user, comprising:
  - holding in memory a function generation unique value  $s$  by a right issuer for the user;
  - creating a value generation unique value  $u$  in a unique value calculation unit from the function generation unique value  $s$  provided from the memory and the unique value  $d$ , the value generation unique value  $u$  being provided as a series of  $m$  values where  $u = (u_1, \dots, u_m)$  to a token for the user;
  - creating by a hash value calculation unit a one-way function value  $X(M)$  of a message  $M$  by applying the one-way function  $H$  to the value generation unique value  $u$  from the unique value calculation unit and the message  $M$ , where the one-way function value  $X(M) = H(u_1 | M) | \dots | H(u_m | M)$ ;
  - holding a certificate  $C$  to prove a public key  $y$  paired with the one-way function value  $X(M)$ ;
  - issuing a capability  $\chi$  from the right issuer to the user, the capability  $\chi$  representing a right of the user in association with the message  $M$ ; and
  - verifying the user from the public key  $y$  and the capability  $\chi$  by a right verifier.
2. (Original) The method for generating a one-way function according to claim 1, wherein the value generation unique value  $u$  is calculated by applying a one-way function  $G$  to the function generation unique value  $s$  and the unique value  $d$ .
3. (Original) The method for generating a one-way function according to claim 1, wherein the value generation unique value  $u$  is calculated by applying an encryption

function E of a symmetric key to the function generation unique value s and the unique value d.

4. (Previously Presented) The method for generating a one-way function according to claim 1, wherein the one-way function value  $X(M)$  of the message M is calculated by applying the one-way function H and an encryption function E of a symmetric key to the value generation unique value u and the message M.

5. (Currently Amended) ~~A device-computer system that generates for generating~~ one-way function values that calculates a one-way function X dependent on a unique value d for a user, comprising:

means for inputting the unique value d;

means for inputting a message M;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d, the value generation unique value u being provided as a series of m values where  $u = (u_1, \dots, u_m)$  to a token for the user;

means for creating a one-way function value  $X(M)$  of the message M by applying a one-way function H to the value generation unique value u from the u-creating means and the message M, where the one-way function value  $X(M) = H(u_1 | M) | \dots | H(u_m | M)$ ;

means for holding a certificate C to prove a public key y paired with the one-way function value  $X(M)$ ;

means for issuing a capability  $\chi$  from the right issuer to the user, the capability  $\chi$  representing a right of the user in association with the message M; and

means for verifying the user from the public key  $y$  and the capability  $\chi$ .

6. (Currently Amended) The ~~device-computer system~~for generating one-way function values according to claim 5, wherein the process of calculating the value generation unique value  $u$  and the one-way function value  $X(M)$  is difficult to observe from the outside.

7. (Currently Amended) A computer system that performs~~proving device for performing~~ processing based on a private key for a user dependent on a message  $M$ , comprising:

means for inputting the message  $M$ , the message  $M$  including at least identifiers of private key processing algorithms;

means for holding a value generation unique value  $u$  for the user;

means for creating a one-way function value  $X(M)$  of the message  $M$  by applying a one-way function  $H$  to the value generation unique value  $u$  from the holding means and the message  $M$ ;

means for performing processing based on the private key and the one-way function value  $X(M)$ ;

means for issuing a capability  $\chi$  from the right issuer to the user, the capability  $\chi$  representing a right of the user in association with the message  $M$ ; and

means for verifying the user from a public key  $y$  and the capability  $\chi$ ,

wherein the value generation unique value  $u$  is created from a function generation unique value  $s$  being held and provided by a right issuer and a unique value  $d$  for the user, the value generation unique value  $u$  being provided as a series of  $m$  values where  $u = (u_1, \dots, u_m)$  to a token for the user, and the one-way function value  $X(M) = H(u_1 \mid M) \mid \dots \mid H(u_m \mid M)$ ; and

wherein the identifiers in the message  $M$  enable the private key processing algorithms to be modified.

8. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the calculation process in processing based on the value generation unique value  $u$  and the one-way function value  $X(M)$  is difficult to observe from the outside.

9. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the ~~proving device~~ computer system is configured as one of a small portable operation device ~~such as and~~ a smart card.

10. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the computer system ~~proving device~~ is configured as a module inside a CPU of a computer ~~the device~~.

11. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the means for performing processing based on the private key comprises:

means for inputting a challenge  $c$ ;

means for calculating a response  $r$  from the challenge  $c$  and the one-way function value  $X(M)$ ; and

means for outputting the response  $r$ .

12. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the means for performing processing based on a private key comprises:

means for inputting a challenge  $c$ ;

means for generating a random number  $k$ ;

means for calculating a response  $r$  from the random number  $k$ , the challenge  $c$ , and the one-way function value  $X(M)$ ; and

means for outputting the response  $r$ .

13. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the means for performing processing based on a private key comprises:

means for generating a random number  $k$ ;

means for calculating a commitment  $w$  from the random number  $k$ ;

means for inputting a challenge  $c$ ;

means for calculating the response  $r$  from the random number  $k$ , the challenge  $c$ , and the one-way function value  $X(M)$ ; and

means for outputting the response  $r$ .

14. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the means for performing processing based on a private key comprises:

means for generating a random number  $k$ ;

means for calculating a commitment  $w$  from the random number  $k$ ;

means for outputting the commitment  $w$ ;

means for inputting a challenge  $c$ ;

means for calculating a response  $r$  from the random number  $k$ , the commitment  $w$ , the challenge  $c$ , and the one-way function value  $X(M)$ ; and

means for outputting the response  $r$ .

15. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the means for performing processing based on a private key performs multiplications and power operations of multiplicative groups on a finite field.

16. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the means for performing processing based on a private key performs additions and scalar multiplication operations of elliptic curves on a finite field.

17. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the means for performing processing based on a private key performs multiplicative residue operations and power residue operations modulo  $n$ , where  $n$  is a composite number that is difficult to factorize.

18. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the message M includes use conditions and the means for inputting messages rejects message input if the use conditions included in the message M are not satisfied.

19. (Currently Amended) The computer system ~~proving device~~ according to claim 7, wherein the message M includes private key processing parameters, and the means for performing processing based on a private key performs processing based on the private key processing parameters included in the message M.

20. (Currently Amended) A ~~device~~ computer system that issues ~~for issuing~~ a proving instrument T in accordance with a unique value d for a user, comprising:

means for inputting the unique value d;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d, the value generation unique value u being provided as a series of m values where  $u = (u_1, \dots, u_m)$  to a token for the user;

means for writing the value generation unique value u from the u-creating means to the proving instrument T;

means for issuing the proving instrument T that includes a hash function X dependent on the unique value d;

means for issuing a capability  $\chi$  from the right issuer to the user, the capability  $\chi$  representing a right of the user in association with the message M; and

means for verifying the user from a public key y and the capability  $\chi$ .

wherein the proving instrument T holds the value generation unique value  $u$ , and upon input of a message  $M$ , creates a one-way function value  $X(M)$  of the message  $M$  by applying a one-way function  $H$  to the value generation unique value  $u$  and the message  $M$  to perform processing based on the one-way function value  $X(M)$  expressed by  $H(u_1 | M) | \dots | H(u_m | M)$ .

21. (Previously Presented) An authentication method by which a right issuer issues rights to right recipients in association with a message  $M$  and a right verifier verifies the rights of the right recipients, the method comprising:

creating a value generation unique value  $u$  from a function generation unique value  $s$  being held and provided by a function generation unique value memory and a unique value  $d$  for a user corresponding to the right recipients, the value generation unique value  $u$  being provided as a series of  $m$  values where  $u = (u_1, \dots, u_m)$  to a token for the user;

calculating a one-way function value  $X(M)$  of the message  $M$  by a hash value generator by applying a one-way function  $H$  to the value generation unique value  $u$  and the message  $M$ , where the one-way function value  $X(M) = H(u_1 | M) | \dots | H(u_m | M)$ ;

issuing a certificate  $C$  to prove a public key  $y$  paired with the one-way function value  $X(M)$  to the right recipients by a certificate issuing unit;

presenting the certificate  $C$  from the right recipients to the right verifier;

performing processing by a private key processing unit based on the one-way function value  $X(M)$ ;

verifying the certificate  $C$  by a certificate verification unit; and

verifying the processing by a private key processing verification unit based on the one-way function value  $X(M)$  of the right recipients with a public key  $y$  proved by the certificate  $C$ .

22. (Original) The authentication method according to claim 21, wherein an identifier aid indicating an authentication type is included in the certificate C issued by the right issuer and the right verifier succeeds in verifying the certificate C only when the authentication identifier aid included in the certificate C matches the type of authentication to be performed.

23. (Original) The authentication method according to claim 21, wherein use conditions are included in the certificate C issued by the right issuer and the right verifier succeeds in verifying the certificate C only when the use conditions included in the certificate C are satisfied.

24. (Currently Amended) A computer system that issues certificate-issuing device for issuing a certificate C in accordance with a unique value d for a user and a message M, comprising:

means for inputting the unique value d;

means for inputting the message M;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d, the value generation unique value u being provided as a series of m values where  $u = (u_1, \dots, u_m)$  to a token for the user;

means for creating a one-way function value  $X(M)$  of the message M by applying a one-way function H to the value generation unique value u from the u-creating means and the message M, where the one-way function value  $X(M) = H(u_1 \parallel M) \parallel \dots \parallel H(u_m \parallel M)$ ;



means for creating a public key  $y$  paired with the one-way function value

$X(M)$ ;

means for issuing a certificate  $C$  to prove the public key  $y$ ;

means for issuing a capability  $\chi$  to the user, the capability  $\chi$  representing a right of the user in association with the message  $M$ ; and

means for verifying the user from the public key  $y$  and the capability  $\chi$ .

25. (Currently Amended) A computer system that performs ~~An authentication device for performing~~ authentication in accordance with a message  $M$ , comprising:

means for inputting the message  $M$ ;

means for holding a value generation unique value  $u$  for a user;

means for creating a one-way function value  $X(M)$  of the message  $M$  by applying a one-way function  $H$  to the value generation unique value  $u$  from the holding means and the message  $M$ ;

means for performing processing based on the one-way function value  $X(M)$ ;

means for holding a certificate  $C$  to prove a public key  $y$  paired with the one-way function value  $X(M)$ ;

means for verifying the certificate  $C$ ;

means for issuing a capability  $\chi$  to the user, the capability  $\chi$  representing a right of the user in association with the message  $M$ ;

means for verifying the user from the public key  $y$  and the capability  $\chi$ ; and

means for verifying processing based on a private key of the user with the public key  $y$ ,

wherein the value generation unique value  $u$  is created from a function generation unique value  $s$  being held and provided by a right issuer and the unique value  $d$  for the user, the value generation unique value  $u$  being provided as a series of  $m$  values where  $u =$

$(u_1, \dots, u_m)$  to a token for the user, and where the one-way function value  $X(M) = H(u_1 \parallel M) \parallel \dots \parallel H(u_m \parallel M)$ .

26. (Previously Presented) An authentication method by which a right issuer issues rights to right recipients in association with a message  $M$  and a right verifier verifies the rights of the right recipients, the method comprising:

creating a value generation unique value  $u$  from a function generation unique value  $s$  being held and provided by a function generation unique value memory and a unique value  $d$  for a user corresponding to the right recipients, the value generation unique value  $u$  being provided as a series of  $m$  values where  $u = (u_1, \dots, u_m)$  to corresponding tokens for the right recipients;

calculating a one-way function value  $X(M)$  of the message  $M$  by a hash value generator by applying a one-way function  $H$  to the value generation unique value  $u$  from the right issuer and the message  $M$ ;

issuing an access ticket  $t$  by an access ticket issuing unit determined from a private key  $x$  and the one-way function value  $X(M)$  to the right recipients, where  $X(M) = H(u_1 \parallel M) \parallel \dots \parallel H(u_m \parallel M)$ ;

performing processing by a private key processing unit based on the one-way function value  $X(M)$ ;

converting the processing by a private key processing conversion unit based on the one-way function value  $X(M)$  to the processing based on the private key  $x$  by the access ticket  $t$ ; and

verifying the processing by a private key processing verification unit based on the one-way function value  $X(M)$  of the right recipients with a public key  $y$  paired with the private key  $x$  by the right verifier.

27. (Original) The authentication method according to claim 21, wherein an identifier aid indicating an authentication type is included in the message M.

28. (Currently Amended) ~~A computer system that issues~~An access ticket issuing device for issuing an access ticket in accordance with a unique value d for a user and a message M, comprising:

means for inputting the unique value d;

means for inputting the message M;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d, the value generation unique value u being provided as a series of m values where  $u = (u_1, \dots, u_m)$  to a token for the user;

means for creating a one-way function value  $X(M)$  of the message M by applying a one-way function H to the value generation unique value u and the message M, where the one-way function value  $X(M) = H(u_1 | M) | \dots | H(u_m | M)$ ;

means for creating the access ticket t from a private key x and the one-way function value  $X(M)$ ;

means for issuing the access ticket t;

means for issuing a capability  $\chi$  from the right issuer to the user, the capability  $\chi$  representing a right of the user in association with the message M; and

means for verifying the user from a public key y and the capability  $\chi$ .

29. (Currently Amended) ~~The access ticket issuing device~~The computer system according to claim 28, wherein the access ticket t is calculated as a difference  $(x - X(M))$  between the private key x and the one-way function value  $X(M)$ .

30. (Currently Amended) The computer system ~~The access ticket issuing device~~ according to claim 28, wherein the access ticket  $t$  is calculated as a quotient  $x/X(M)$  between the private key  $x$  and the one-way function value  $X(M)$ .

31. (Currently Amended) The computer system ~~The access ticket generation device~~ according to claim 28, wherein the unique value  $d$  for the user is  $(d_1, \dots, d_m)$ , the value generation unique value  $u$  is  $(u_1, \dots, u_m)$  and the one-way function value  $X(M)$  is generated from bit concatenation  $H(u_1 \parallel M) \parallel \dots \parallel H(u_m \parallel M)$  of the value of the one-way function  $H$  and has a desired bit length.

32. (Currently Amended) The computer system ~~The access ticket generation device~~ according to claim 31, wherein the value generation unique value  $(u_1, \dots, u_m)$  is found from  $u_j = G(s_j \parallel d)$  obtained by applying a one-way function  $G$  to the function generation unique value  $s = (s_1, \dots, s_m)$ .

33. (Currently Amended) A computer system that performs ~~An authentication device for performing~~ authentication for a user in accordance with a message  $M$ , comprising:

means for inputting the message  $M$ ;

means for holding a value generation unique value  $u$  for the user;

means for creating a one-way function value  $X(M)$  of the message  $M$  by applying a one-way function  $H$  to the value generation unique value  $u$  from the holding means and the message  $M$ ;

means for performing processing based on the one-way function value  $X(M)$ ;

means for holding an access ticket  $t$  determined from a private key  $x$  and the one-way function value  $X(M)$ ;

means for converting the processing based on the one-way function value  $X(M)$  to processing based on the private key  $x$  by the access ticket  $t$ ;

means for holding a public key  $y$  paired with the private key  $x$ ;

means for issuing a capability  $\chi$  from the right issuer to the user, the capability  $\chi$  representing a right of the user in association with the message M;

means for verifying the user from the public key y and the capability  $\chi$ ; and

means for verifying the processing based on the private key x with the public key y,

wherein the value generation unique value u is created from a function generation unique value s being held and provided by a right issuer and a unique value d provided for the user, the value generation unique value u being provided as a series of m values where  $u = (u_1, \dots, u_m)$  to a token for the user, and where the one-way function value  $X(M) = H(u_1 | M) | \dots | H(u_m | M)$ .

34. (Currently Amended) The authentication device~~computer system~~ according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a challenge c with the access ticket t.

35. (Currently Amended) The computer system~~authentication device~~ according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a response r with the access ticket t.

36. (Currently Amended) The computer system~~authentication device~~ according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a response r with the access ticket t and a challenge c.

37. (Currently Amended) The computer system~~authentication device~~ according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a challenge c with a commitment w and means for updating a response r with the access ticket t and the challenge c.

38. (Currently Amended) The computer system~~authentication device~~ according to claim 33, wherein the means for converting the processing based on the private key comprises

means for updating a challenge  $c$  with the access ticket  $t$  and a commitment  $w$ , and means for updating a response  $r$  with the commitment  $w$ .